



Approach for Continuous and Transparent User Identity Verification for Secure Internet Services

U RAVI¹, P.NAGESWARARAO²

¹M.Tech Student, Sree Rama institute of technology and science

Kuppenakuntla, Penuballi, Khammam, TS INDIA

²Asst Prof, CSE Dept Sree Rama institute of technology and science

Kuppenakuntla, Penuballi, Khammam, TS INDIA

ABSTRACT:

Security of the web based services is become serious concern now a days. Secure user authentication is very important and fundamental in most of the systems User authentication systems are traditionally based on pairs of username and password and verify the identity of the user only at login phase. No checks are performed during working sessions, which are terminated by an explicit logout or expire after an idle activity period of the user. Emerging biometric solutions provides substituting username and password with biometric data during session establishment, but in such an approach still a single shot verification is less sufficient, and the identity of a user is considered permanent during the entire session. A basic solution is to use very short session timeouts and periodically request the user to input his credentials over and over, but this is not a definitive solution and heavily penalizes the service usability and ultimately the satisfaction of users. This paper explores promising alternatives offered by applying biometrics in the management of sessions. A secure protocol is defined for perpetual authentication through continuous user verification. Finally,

the use of biometric authentication allows credentials to be acquired transparently i.e. without explicitly notifying the user or requiring his interaction, which is essential to guarantee better service usability.

Index Terms—Cloud storage, data sharing, key-aggregate encryption, patient-controlled encryption.

INTRODUCTION:

CLOUD storage is gaining popularity recently. In enterprise settings, we see the rise in demand for data outsourcing, which assists in the strategic management of corporate data. It is also used as a core technology behind many online services for personal applications. Nowadays, it is easy to apply for free accounts for email, photo album, file sharing and/or remote access, with storage size more than 25 GB (or a few dollars for more than 1 TB). Together with the current wireless technology, users can access almost all of their files and emails by a mobile phone in any corner of the world. Considering data privacy, a traditional way to ensure it is to rely on the server to enforce the access control after authentication (e.g., [1]), which means any unexpected privilege escalation will expose all data. In a shared-tenancy cloud computing environment, things become even worse. Data from different clients can be hosted on separate virtual machines (VMs) but reside on a single physical machine. Data in a target VM could be stolen by instantiating another VM coresident with the target one. Regarding availability of files, there are a series of cryptographic schemes which go as far as allowing a third-party auditor to check the availability of files on behalf of the data owner without leaking anything about the data [3], or without compromising the data owners anonymity. Likewise, cloud users probably will not hold the strong belief that the cloud server is doing a good job in terms of confidentiality. A cryptographic solution, for example, [5], with proven security relied on number theoretic assumptions is more desirable, whenever the user is not perfectly happy with trusting the security of the VM or the honesty of the technical staff. These

users are motivated to encrypt their data with their own keys before uploading them to the server.

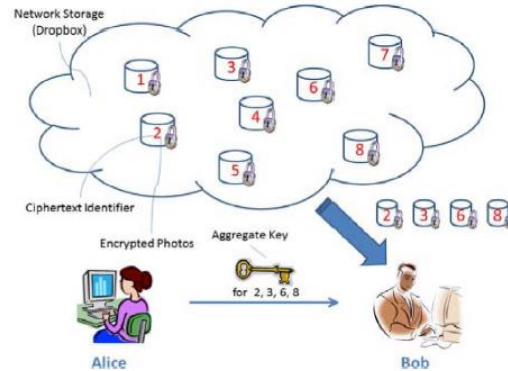


Fig. 1. Alice shares files with identifiers 2, 3, 6, and 8 with Bob by sending him a single aggregate key.

Existing System

We first give the framework and definition for key aggregate encryption. Then we describe how to use KAC in a scenario of its application in cloud storage. A key-aggregate encryption scheme consists of five polynomial-time algorithms as follows. The data owner establishes the public system parameter via Setup and generates a public/master-secret key pair via KeyGen. Messages can be encrypted via Encrypt by anyone who also decides what ciphertext class is associated with the plaintext message to be encrypted. The data owner can use the master-secret to generate an aggregate decryption key for a set of ciphertext classes via Extract. The generated keys can be passed to delegates securely (via secure e-mails or secure devices) Finally, any user with an aggregate key can decrypt any ciphertext provided that the ciphertext's class is contained in the aggregate key via Decrypt.

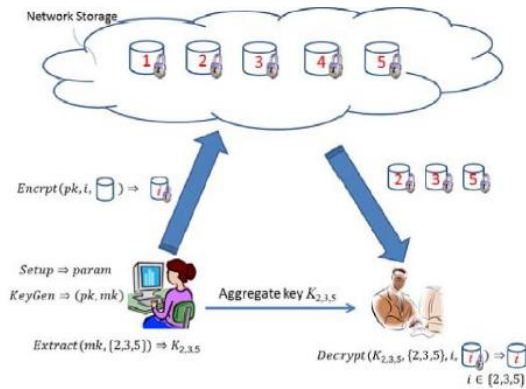


Fig. 2. Using KAC for data sharing in cloud storage.

Proposed System:

This section we compare our basic KAC scheme with other possible solutions on sharing in secure cloud storage. We summarize our comparisons in Table We take the tree structure as an example. Alice can first classify the ciphertext classes according to their subjects like Fig. 3. Each node in the tree represents a secret key, while the leaf nodes represents the keys for individual ciphertext classes. Filled circles represent the keys for the classes to be delegated and circles circumented by dotted lines represent the keys to be granted. Note that every key of the non leaf node can derive the keys of its descendant nodes. In Fig. 3a, if Alice wants to share all the files in the “personal” category, she only needs to grant the key for the node “personal,” which automatically grants the delegate the keys of all the descendant nodes (“photo,” “music”). This is the ideal case, where most classes to be shared belong to the same branch and thus a parent key of them is sufficient.

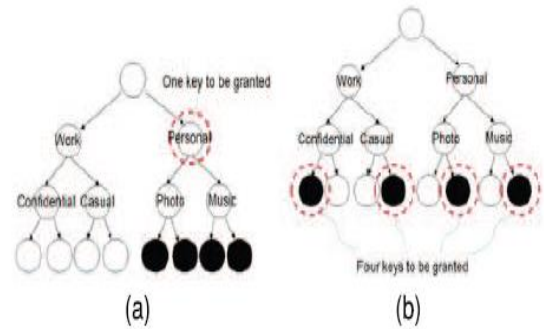
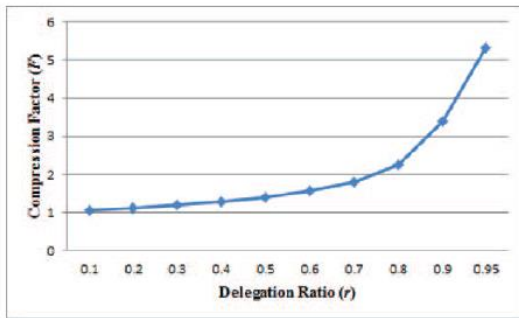


Fig. 3. Compact key is not always possible for a fixed hierarchy.

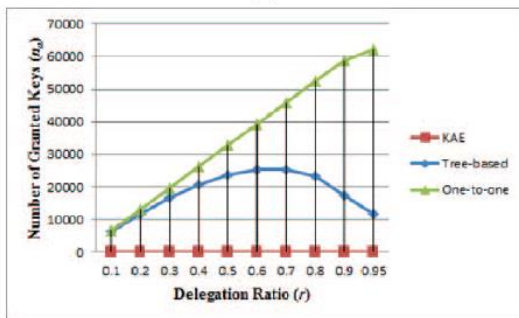
PERFORMANCE ANALYSIS

Compression Factors

For a concrete comparison, we investigate the space requirements of the tree-based key assignment approach we described in Section 3.1. This is used in the complete subtree scheme, which is a representative solution to the broadcast encryption problem following the well-known subset-cover framework [33]. It employs a static logical key hierarchy, which is materialized with a full binary key tree of height h (equals to 3 in Fig. 3), and thus can support up to $2h$ ciphertext classes, a selected part of which is intended for an authorized delegatee. In an ideal case as depicted in Fig. 3a, the delegatee can be granted the access to $2h_s$ classes with the possession of only one key, where h_s is the height of a certain subtree (e.g., $h_s = 2$ in Fig. 3a). On the other hand, to decrypt ciphertexts of a set of classes, sometimes the delegatee may have to hold a large number of keys, as depicted in Fig. 3b. Therefore, we are interested in n_a , the number of symmetrickeys to be assigned in this hierarchical key approach, in an average sense.



(a)



(b)

Fig. 5. (a) Compression achieved by the tree-based approach for delegating different ratio of the classes.
(b) Number of granted keys (n_a) required for different approaches in the case of 65,536 classes of data.

NEW PATIENT-CONTROLLED ENCRYPTION (PCE)

Motivated by the nationwide effort to computerize America's medical records, the concept of patient-controlled encryption has been studied [8]. In PCE, the health record is decomposed into a hierarchical representation based on the use of different ontologies, and patients are the parties who generate and store secret keys. When there is a need for a healthcare personnel to access part of the record, a patient will release the secret key for the concerned part of the record. In the work of Benaloh et al. [8], three solutions have been provided, which are symmetric-key PCE for fixed hierarchy (the "folklore" tree-based method in Section 3.1), public-

key PCE for fixed hierarchy (the IBE analog of the folklore method, as mentioned in Section 3.1), and RSAbased symmetric-key PCE for "flexible hierarchy" (which is the "set membership" access policy as we explained).

CONCLUSION

This paper provides various existing methods used for continuous authentication using different biometrics. Initial one time login verification is inadequate to address the risk involved in post logged in session. Therefore this paper attempts to provide a comprehensive survey of research on the underlying building blocks required to build a continuous biometric authentication system by choosing bio-metric. Continuous authentication verification with multi-modal biometrics improves security and usability of user session

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U RAVI is an M.Tech Department of Computer Science & Engineering, Sreerama Institute of Technology & science, Penuballi Mandal, Khammam, Kotha Kuppenkuntla



Mr. P.Nageswa Rao is an efficient teacher, received M.Tech from JNTU Hyderabad is working as an Assistant Professor in Department of C.S.E, Sree Rama Institute of Technology & Science,

Kuppenakuntla, Penuballi, Khammam, AP,India. He has published many papers in both National & International Journals. His area of Interest includes Data Communications & Networks, Database Management Systems, Computer Organization, C Programming and other advances in Computer Applications